

# Internship Report – M2 MPRI

## Mnemonic Monads and Compactness

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## General context

We work in the setting of 2-category theory to understand equivalences given by taking *prime* or *compact* elements appearing in denotational semantics and categorical logic. These appear under the name of compact/finite/finitely presented object, atomic object, join irreducible, join prime. We will show that these can all be put into a common picture using the language of lax-idempotent (KZ) monads [Koc95; Zöb76], which are commonly thought of as *free cocompletions* under a class of colimits. In previous work these ideas have been considered in isolation. A first step towards this can be seen in [DL18] where the authors recover duality theorems by axiomatizing accessibility.

## Research problem

This internship was aimed at introducing the study of *mnemonic lax-idempotent monads*, which provide a particularly pleasant contexts for studying cocompletions. The specific question was to find out when the inclusion from the basic data to the continuous algebras of the monad is a local equivalence. This question naturally set the stage for the general study of what a generator is with respect to a lax-idempotent monad.

This is a first step towards a general theory of generators for lax-idempotent monads, abstracting and unifying several constructions previously treated separately, such as in domain theory, realizability, and categorical logic. In a similar vein to [ADL23; DL18], we develop formal category theory in the setting of lax-idempotent pseudomonads.

Using this approach we obtain a synthetic approach to studying generators, that is without regard to the base 2-category. The main application so far is to recover results showing that certain objects are freely generated via different cocompletions. We also recover new notions of generator in the setting of free opfibrations on a functor [Koc13].

## Contribution

We generalized the ideas of compact element, way below relation and having enough compact elements from domain Theory to the setting of lax-idempotent monads. These constructions are new and moreover give a conceptual answer to the original questions, namely taking generators is right adjoint to taking free continuous algebras. The advantage of the gained generality is that we can now capture many standard examples using the same technology while being able to formally prove theorems about them. The approach was to have a collection of examples from a wide range of topics from domain theory and topos theory to fibrations and realizability. This way we were able to find the common denominator of the constructions.

## Arguments supporting its validity

The work has been mathematical in nature, thus all the claims come with proofs or references to their proofs. We claim that the given definitions are appropriate because of:

1. the generality and simplicity of the proofs they give rise to
2. the wide range of applications.

These applications range from lattice theory all the way to topos theory, type theory and higher category theory.

Our results rely on the observation that in **CAT** and in **Pos**, Kan extensions are computed pointwise, this assumption breaks down in a lot of other 2-Categories. Nonetheless the examples of **CAT** and **Pos** contain all the classical results we wanted to generalize and we plan on making the same treatment work for other 2-Categories by considering Virtual Equipments [CS09].

## Summary and future work

We used the technology of lax-idempotent pseudomonads on 2-categories to define what it means to be a generator for an algebra using purely arrow theoretic arguments. We call this notion T-compactness, that is a notion of generator indexed by a given lax-idempotent pseudomonad. In the case of the **Ind**-construction, T-compactness recovers standard notion of compact/finitely presented object. Using this we developed the tools to study new notions of generator using a simple recipe.

More surprisingly we found that our treatment of lax-idempotent pseudomonads is a kind of lax-idempotent generalization of the work on monadic descent by [Mes06].

The main contribution is theorem 5.15, stating that taking generators is right adjoint to taking free continuous algebras. This gives a conceptual way of understanding equivalences between cauchy completions and T-compactly generated objects.

The next step will be to work in the more appropriate setting of virtual equipments to capture more exotic examples of generator. We have identified two such examples that we expect to yield new theorems with applications to semantics of dependent type theory, realizability and higher category theory. Namely the Monoidal Categories are lax-idempotent over Multicategories [Her01], and Opfibrations are lax-idempotent over **CAT/B** [Koc13].

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# 1 Introduction

A *cocompletion* of a category is a universal way of freely adjoining certain colimits. Under mild size conditions each of these cocompletions give rise to a pseudomonad on locally small categories. These satisfy a certain adjointness property which makes them more pleasant to work with, these are known as *lax-idempotent (or KZ) pseudomonads* [Koc95; Zöb76]. This lax-idempotent pseudomonad makes sense for other 2-categories than just locally small categories, for example partially ordered sets admit many naturally occurring lax-idempotent pseudomonads. Among them is the *sup-cocompletion* adjoining arbitrary joins to a poset, this is done using the *down-set construction*. The goal of this report is to develop a general framework for understanding how cocompletions give rise to “generators” through their associated lax-idempotent pseudomonads. For example in the case of the *directed-join-cocompletion* on posets, we obtain the standard notions of *finite element* and *way below relation* from domain theory.

## 2 Background

### 2.1 2-Categories

A 2-category is morally a category with arrows between the arrows, the prime example being CAT which has objects locally small categories, morphisms functors and 2-cells natural transformations.

**Definition 2.1.** A 2-category  $\mathcal{B}$  consists of:

- a class of objects (denoted  $A, B, C$ ),
- for each pair  $A, B$ , a category  $\mathcal{B}(A, B)$ , whose objects are called **arrows**, and morphisms **2-cells**. We the composition of 2-cells is called **vertical composition**,
- for each object  $A$  an **identity arrow**  $\text{id}_A \in \mathcal{B}(A, A)$ ,
- for each triple of objects  $A, B, C$  a **horizontal composition** functor  $(- \circ_{A,B,C} -) : \mathcal{B}(B, C) \times \mathcal{B}(A, B) \rightarrow \mathcal{B}(A, C)$ ,

making the following diagrams commute:

1. Associativity:

$$\begin{array}{ccc}
 \mathcal{B}(C, D) \times \mathcal{B}(B, C) \times \mathcal{B}(A, B) & \xrightarrow{\text{id}_{\mathcal{B}(C,D)} \times \circ_{A,B,C}} & \mathcal{B}(C, D) \times \mathcal{B}(A, C) \\
 \downarrow \circ_{B,C,D} \times \text{id}_{\mathcal{B}(A,B)} & & \downarrow \circ_{A,C,D} \\
 \mathcal{B}(B, D) \times \mathcal{B}(A, B) & \xrightarrow{\circ_{A,B,D}} & \mathcal{B}(A, D)
 \end{array}$$

2. Unit Laws:

$$\begin{array}{ccc}
 \mathcal{B}(A, B) & \xrightarrow{(\text{id}_{\mathcal{B}(A,B)}, \text{id}_A)} & \mathcal{B}(A, B) \times \mathcal{B}(A, A) \\
 \downarrow (\text{id}_B, \text{id}_{\mathcal{B}(A,B)}) & \searrow \text{id}_{\mathcal{B}(A,B)} & \downarrow \circ_{A,A,B} \\
 \mathcal{B}(B, B) \times \mathcal{B}(A, B) & \xrightarrow{\circ_{A,B,B}} & \mathcal{B}(A, B)
 \end{array}$$

We will always write  $f \circ g$  instead of  $f \circ_{A,B,C} g$ . For 2-cells we write vertical composition as  $\alpha \circ \beta$  and horizontal composition as  $\alpha \star \beta$ . We also use the standard notation for whiskering  $\alpha f := \alpha \star \text{id}_f$  and  $f \alpha := \text{id}_f \star \alpha$  for an arrow  $f$  and a 2-cell  $\alpha$  which are compatible. Horizontal composition has lower priority than vertical composition, i.e. is short for

$$\alpha \star \beta \circ \gamma \star \delta = (\alpha \star \beta) \circ (\gamma \star \delta).$$

**Example 2.2.** The 2-category **Pos** has objects partially ordered sets  $(X, \leq)$  and hom-categories  $\text{Pos}(X, Y)$  non-decreasing functions with their pointwise order viewed as a category.

**Example 2.3.** The 2-category **CAT** has as objects locally small categories and functors as arrows and natural transformations as 2-cells. Similarly 2-category **Lex** has as objects finitely complete locally small categories and finite limit preserving functors as arrows and natural transformations as 2-cells.

**Example 2.4.** For any 2-category  $\mathcal{B}$  given an object  $A \in \mathcal{B}_0$ , we can construct the **slice-2-category**  $\mathcal{B}/A$  which has objects arrows with codomain  $A$ , morphisms commuting triangles and 2-cells are of the form

$$\begin{array}{ccc} X & \xrightarrow{f} & Y \\ & \Downarrow \alpha & \\ X & \xrightarrow{g} & Y \\ a \swarrow & & \searrow b \\ & A & \end{array},$$

i.e. 2-cells in  $\mathcal{B}$  such that  $b\alpha = a$  for  $a, b \in \mathcal{B}/A$ .

## 2.2 Universal Constructions in a 2-Category

**Definition 2.5.** An **adjunction** in a 2-category  $\mathcal{B}$  is given by a pair of arrows  $l : A \rightarrow B$  and  $r : B \rightarrow A$  together with 2-cells  $\eta : \text{id}_A \Rightarrow rl$  and  $\epsilon : lr \Rightarrow \text{id}_B$  such that

$$\epsilon l \circ l \eta = \text{id}_l$$

and

$$r \epsilon \circ \eta r = \text{id}_r.$$

We denote this situation by  $l \dashv r$ .

**Definition 2.6.** [SW78] Given arrows  $f : A \rightarrow X$  and  $k : A \rightarrow B$  in a 2-category  $\mathcal{B}$ , a **left extension of  $f$  along  $k$**  is a universal pair  $(l : B \rightarrow X, \alpha : f \Rightarrow lk)$ : That is for any other pair  $h : B \rightarrow X, \beta : f \Rightarrow hk$  there is a unique  $\gamma : l \Rightarrow h$  such that  $\beta = \gamma k \circ \alpha$ .

$$\begin{array}{ccc} \begin{array}{ccc} B & & \\ \uparrow k & \nearrow \beta & \searrow h \\ A & \xrightarrow{f} & X \end{array} & = & \begin{array}{ccc} B & & \\ \uparrow k & \nearrow \alpha & \searrow h \\ A & \xrightarrow{f} & X \end{array} \end{array}$$

We write  $\text{lan}_k f$  for the arrow of the left extension of  $f$  along  $p$  when it exists.

A left extension  $(l, \alpha)$  of  $f$  along  $k$  is **absolute** if given any arrow  $g : X \rightarrow Y$ ,  $(gl, g\alpha)$  is a left extension of  $gf$  along  $k$ .

The following are the lesser known duals of extensions, a good reference is [Di +17].

**Definition 2.7.** [SW78] Given arrows  $f : A \rightarrow B$  and  $p : E \rightarrow B$  in a 2-category  $\mathcal{B}$ , a **left lift of  $f$  along  $p$**  is a universal pair  $(l : A \rightarrow E, \alpha : f \Rightarrow pl)$ : That is for any other pair  $h : A \rightarrow E, \beta : f \Rightarrow ph$  there is a unique factorization  $\gamma : l \Rightarrow h$  such that  $\beta = p\gamma \circ \alpha$ .

We write  $\text{lift}_p f$  for the arrow of the left lift of  $f$  along  $p$  when it exists.

A left lift  $(l, \alpha)$  of  $f$  along  $p$  is **absolute** if given any arrow  $g : C \rightarrow A$ ,  $(lg, \alpha g)$  is a left lift of  $fg$  along  $p$ .

The definition is exactly saying that we have a left extension in the opposite 2-category, thus left lifts are unique up to unique iso. Of course we can dualize all these to obtain right extensions and right lifts in the same way.

The following lemma states that adjunctions are a special case of left lifts.

**Lemma 2.8.** (2.4 of [Di +17]) Given arrows  $l : A \rightarrow B$  and  $r : B \rightarrow A$  such that  $l \dashv r$  in  $\mathcal{B}$  with unit  $\eta : id_A \Rightarrow rl$ , then  $(l, \eta)$  is an absolute left lift of  $r$  along  $id_A$ . Conversely if  $(\eta, l)$  is an absolute left lift of  $r$  along  $id_A$ , then we can construct a counit  $\epsilon : lr \Rightarrow id_B$  such that  $(l, r, \eta, \epsilon)$  is an adjunction.

We will make extensive use of inverters which should be thought of as a 2-categorical analogue of an equalizer which can take 2-cells into account. A more thorough treatment of inverters can be found in Section B1.15 of [Joh02].

**Definition 2.9.** In a 2-category  $\mathcal{B}$  with arrows  $f, g : A \rightarrow B$  and a 2-cell  $\alpha : f \Rightarrow g$ , an **inverter** of  $\alpha$  is an arrow  $\iota : X \rightarrow A$  such that:

- $\alpha\iota$  is invertible.
- For every object  $Y \in \mathcal{B}$ , the functor  $- \circ \iota : \mathcal{B}(Y, X) \rightarrow \mathcal{B}(Y, A)$  is full and faithful with its replete image in arrows which invert  $\alpha$ , that is if  $\alpha f$  is invertible for some  $f : Y \rightarrow A$ , then there is some  $\tilde{f} : X \rightarrow A$  such that  $\iota \circ \tilde{f} \cong f$ .

### 3 Basics of Lax-Idempotent Monads

#### 3.1 Lax-Idempotent Monads as a Setting for Cocompletions

Lax-idempotent pseudomonads capture the idea of *cocompletion*. The most common examples are the down-set completion and the ideal completion on posets and the presheaf construction or the Ind-completion on locally small categories. The organizing power of lax-idempotent monads is for example being used by [Blo25] to generalize common 1-categorical arguments to that of  $(\infty, 1)$ -categories in a painless way.

It can also be seen from [ADL23] that lax-idempotent monads can give a satisfactory context for formal category theory. Moreover we will see that certain dualities can also be viewed

from this perspective.

A definition of pseudomonads can be found in Appendix A.

**Definition 3.1.** A pseudomonad  $(T : \mathcal{B} \rightarrow \mathcal{B}, \mu, \eta)$  on a 2-category  $\mathcal{B}$  is lax-idempotent (or KZ) [Koc95; Zöb76] when for every object  $X \in \mathcal{B}$ , we have an adjunction

$$\mu_X \dashv \eta_{TX},$$

such that the invertible counit is given by the left unitor  $l_X : \mu_X \circ \eta_{TX} \Rightarrow \text{id}_{TX}$ . In this case it was shown (cf. [Joh02]), we also get

$$T\eta_X \dashv \mu_X.$$

This situation can then be summarized by saying for each object  $X$ , we obtain two adjunctions

$$\begin{array}{ccc} & T^2X & \\ T\eta_X \swarrow & \downarrow \mu_X & \searrow \eta_{TX} \\ & TX & \end{array}.$$

We will refer to lax-idempotent pseudomonads as **lax-idempotent monads**.

In [MW12], the authors showed that a lax-idempotent monad can be specified very efficiently.

**Definition 3.2.** A left Kan pseudomonad  $T$  on a 2-category  $\mathcal{B}$  is given by

- a function  $T_0 : \mathcal{B}_0 \rightarrow \mathcal{B}_0$
- for each object  $X \in \mathcal{B}$ , an arrow  $\eta_X : X \rightarrow TX$
- for every arrow  $f : X \rightarrow TY$  a left extension along  $\eta_X$ :

$$\begin{array}{ccc} X & \xrightarrow{f} & TY \\ \eta_X \downarrow & \swarrow \text{\scriptsize } T_f \quad \searrow \text{\scriptsize } f^T & \\ TX & & \end{array}$$

satisfying that

1. the following is a left extension for each  $X \in \mathcal{B}$ :

$$\begin{array}{ccc} X & \xrightarrow{\eta_X} & TX \\ \eta_X \downarrow & \swarrow \text{\scriptsize } \text{id}_{\eta_X} \quad \searrow \text{\scriptsize } \text{id}_{TX} & \\ TX & & \end{array}$$



2. for each  $f : X \rightarrow TY$  and  $g : Y \rightarrow TZ$ :

$$\begin{array}{ccccc}
 X & \xrightarrow{f} & TY & \xrightarrow{g^T} & TZ \\
 \eta_X \downarrow & \swarrow T & \nearrow f^T & & \\
 & TX & & & 
 \end{array}$$

, the arrow  $g^T \circ f^T$  is a left extension of  $g^T \circ f$  along  $\eta_X$ .

**Lemma 3.3.** [MW12] Every lax-idempotent monad  $T : \mathcal{B} \rightarrow \mathcal{B}$  arises from a left Kan pseudomonad with the action on arrows given by

$$T(f : A \rightarrow B) := \text{lan}_{\eta_A}(\eta_B f)$$

and the multiplication is given by

$$\mu_X := \text{lan}_{\eta_{TX}} id_{TX}.$$

### 3.2 Many Examples of Lax Idempotent Monads

The following example is the general shape that we will consider here.

**Proposition 3.4.** There is a lax-idempotent monad  $\text{Ind} : \text{CAT} \rightarrow \text{CAT}$  specified by the following left Kan pseudomonad on  $\text{CAT}$ :

- On objects, define  $\text{Ind}(C) \subseteq \text{Set}^{C^{op}}$  defined on the presheaves which are small filtered colimits of representables
- The unit maps  $\eta_C : C \rightarrow \text{Ind}(C)$  are given by the Yoneda embedding restricted on its codomain

Section 3.1 shows that the action on morphisms (functors)  $F : A \rightarrow B$  is given by the following left Kan extension:

$$\begin{array}{ccc}
 A & \xrightarrow{F} & B \\
 \eta_A \downarrow & & \downarrow \eta_B \\
 \text{Ind}(A) & \xrightarrow{\text{lan}_{\eta_A} \eta_B F} & \text{Ind}(B)
 \end{array}$$

Moreover the multiplication is given on objects  $C \in \text{CAT}$  by

$$\mu_C(\psi \in \text{Ind}^2(C)) := \text{colim}(\int \psi \rightarrow \text{Ind}(C)).$$

Here  $\int \psi$  is the category of elements of  $\psi$  [Awo06].

**Example 3.5.** [PT21] There is a lax-idempotent monad  $\mathcal{P} : \text{CAT} \rightarrow \text{CAT}$  specified by the following left Kan pseudomonad on  $\text{CAT}$ :

- On objects, define  $\mathcal{P}(C) \subseteq \text{Set}^{C^{op}}$  defined on the presheaves which are small colimits of representables

- The unit maps  $\eta_C : C \rightarrow \mathcal{P}(C)$  are given by the Yoneda embedding on its codomain

**Example 3.6.** [Gie+80] A subset  $S$  of a poset  $X$  is **down-closed** if  $s \in S$  and  $t \leq s$ , implies  $t \in S$ . A subset  $S$  of a poset  $X$  is **directed** if it is non-empty and for any  $s, t \in S$  there is  $u \in S$  such that  $s \leq u$  and  $t \leq u$ .

There is a lax-idempotent monad  $\text{Idl} : \text{Pos} \rightarrow \text{Pos}$  specified by the following left Kan pseudomonad on  $\text{Pos}$ :

- On objects, define  $\text{Idl}(X) := \{S \subseteq X \mid \text{directed and down-closed}\}$
- The unit maps  $\eta_X : X \rightarrow \text{Idl}(X)$  are given by principal down-sets  $\eta_X(x) := \{y \in X \mid y \leq x\}$ .

### 3.3 The 2-category of Algebras for a Lax Idempotent Monad

For pseudomonads there are multiple notions of algebra, we will only consider the pseudo-algebras (cf. Section A) and content ourselves with the fact that for lax-idempotent monads they are easy to detect. When it is clear from context we will refer to pseudo-T-algebras for a given lax-idempotent monad  $T$  as **algebras**.

**Proposition 3.7.** [Koc95; Zöb76] *The algebras of a lax-idempotent monad  $(T : \mathcal{B} \rightarrow \mathcal{B}, \mu, \eta)$  are objects  $X \in \mathcal{B}$  such that the unit  $\eta_X$  has a left adjoint with invertible counit  $\epsilon$ , that is  $\alpha \dashv \eta_X$  and  $\epsilon : \alpha \eta_X \cong \text{id}$ .*

*Remark 3.8.* This means that being an algebra for a lax-idempotent monad is a **property**.

*Remark 3.9.* Let  $(X, \alpha), (Y, \beta)$  be algebras for a lax-idempotent monad  $T : \mathcal{B} \rightarrow \mathcal{B}$  and let  $f : X \rightarrow Y$  be an arrow  $\mathcal{B}$ . We can always construct a 2-cell  $\tilde{f} : \beta T f \Rightarrow f \alpha$  as follows:

$$\begin{array}{ccccc}
 & & TX & \xrightarrow{Tf} & TY & \xrightarrow{\beta} & Y \\
 & \text{id}_{TX} \curvearrowright & \uparrow \eta_X & \searrow \eta_f & \uparrow \eta_Y & \text{counit} \curvearrowright & \uparrow \text{id}_Y \\
 TX & \xrightarrow{\alpha} & X & \xrightarrow{f} & Y & & 
 \end{array}$$

This is exactly the **mate** [KS06] of  $\eta_f$ .

**Example 3.10.** In the case of  $\text{Idl} : \text{Pos} \rightarrow \text{Pos}$ , take  $X, Y \in \text{Dcpo}$  and consider a non-decreasing function  $f : X \rightarrow Y$ , the 2-cell  $\tilde{f}$  tells us that for any directed set  $S \subset X$ ,  $\bigvee f(S) \leq f \bigvee S$ .

**Definition 3.11.** [KL97] A morphism  $f : (X, \alpha) \rightarrow (Y, \beta)$  of algebras for a lax-idempotent monad is an arrow  $f : A \rightarrow B$  such that the mate  $\tilde{f}$  is invertible.

**Example 3.12.** For any morphism  $f : X \rightarrow Y$ ,  $Tf : TX \rightarrow TY$  is an algebra morphism.

**Proposition 3.13.** [KL97] *For any lax-idempotent monad  $T : \mathcal{B} \rightarrow \mathcal{B}$ , there is a 2-category  $T\text{-Alg}$  which has as objects algebras, morphisms algebra morphisms and as 2-cells the original 2-cells of  $\mathcal{B}$ .*

**Example 3.14.** The monad  $\mathcal{P}$  from 3.5 has as algebras locally small cocomplete categories  $\mathbf{CoCompCat}$ .

The monad  $\mathbf{Ind}$  from 3.6 has as algebras posets admitting all directed joins, these are called **directed complete partial ordered sets (DCPO)**, we denote this 2-category of algebras by  $\mathbf{Dcpo}$ .

The following lemma is a formal version of the fact that left adjoints preserve colimits.

**Proposition 3.15.** *For lax-idempotent monads, left adjoints are always morphisms of algebras.*

*Proof.* Consider algebras  $(X, \alpha)$  and  $(Y, \beta)$  for a lax-idempotent monad  $T : \mathcal{B} \rightarrow \mathcal{B}$  and assume we have a morphism  $l : X \rightarrow Y \in \mathcal{B}$  with a right adjoint  $r : Y \rightarrow X$ . We wish to show that  $l\alpha \cong \beta Tl$  via a canonical isomorphism. To do this notice that  $\alpha \dashv \eta_X$  and  $l \dashv r$ . Now since adjoints compose, we have that

$$l\alpha \dashv \eta_X r$$

And we have that  $\eta_X r \cong \text{Tr}\eta_Y$  by naturality, giving us that

$$\beta Tl \dashv \text{Tr}\eta_Y.$$

Now left adjoints are unique up to unique isomorphism, so  $l\alpha \cong \beta Tl$ . □

**Corollary 3.16.** *For any algebra  $(X, \alpha)$ ,  $\alpha : TX \rightarrow X$  is an algebra morphism.*

## 4 Compactness

Now we introduce a notion of compactness parametrized by a lax-idempotent monad. The idea is to recover the usual notions of generator appearing in the literature, for example projective objects, finitely presented objects or connected objects and unify their study.

**Definition 4.1.** A generalized element  $x : A \rightarrow X$  of an algebra  $(X, \alpha)$  is **T-compact** if  $(\eta_X x, c_X^{-1} : x \Rightarrow \alpha \eta_X x)$  is an absolute left lift of  $x$  along  $\alpha$

$$\begin{array}{ccc} & & TX \\ & \nearrow \eta_X x & \downarrow \alpha \\ A & \xrightarrow{x} & X \end{array}$$

(The diagram shows a triangle with vertices  $A$ ,  $X$ , and  $TX$ . An arrow labeled  $x$  goes from  $A$  to  $X$ . An arrow labeled  $\eta_X x$  goes from  $A$  to  $TX$ . An arrow labeled  $\alpha$  goes from  $TX$  to  $X$ . A double arrow (counit) goes from  $\eta_X x$  to  $\alpha$ .)

where  $c_X : \alpha \eta_X \Rightarrow \text{id}_X$  is the invertible counit of  $\alpha \dashv \eta_X$ . In particular this means that we have a natural bijection we have a natural bijection:

$$\frac{\eta_X x \Rightarrow u}{x \Rightarrow \alpha u}$$

stable under precomposition.

*Remark 4.2.* The way to think about this is to say that around  $x$ ,  $\alpha$  behaves like a right adjoint to  $\eta_X$ .

**Proposition 4.3.** *For the lnd-completion, a finitely presented object  $X \in \mathcal{A}^1$  is exactly one for which  $X : 1 \rightarrow A$  is lnd-compact.*

*Proof.* Assume  $X : 1 \rightarrow \mathcal{A}$  is lnd-compact and consider a diagram  $D : I \rightarrow \mathcal{A}$  with  $I$  a filtered category. Denote  $\psi := \text{colim}(\eta_{\mathcal{A}} \circ D)$  the associated object in  $\text{Ind}(\mathcal{A})$ . So now:

$$\begin{aligned} \mathcal{A}(X, \text{colim} D) &= \mathcal{A}(X, \alpha\psi) \\ &\cong \text{Ind}(\mathcal{A})(\eta_{\mathcal{A}}X, \psi) \\ &\cong \text{Ind}(\mathcal{A})(\eta_{\mathcal{A}}X, \text{colim} \int \psi \rightarrow \mathcal{A} \rightarrow \text{Ind}(\mathcal{A})) \\ &\cong \text{colim}_{i \in I} \text{Ind}(\mathcal{A})(\eta_{\mathcal{A}}X, \eta_{\mathcal{A}}C_i) \\ &\cong \text{colim}_{i \in I} \mathcal{A}(X, C_i) \end{aligned}$$

The second to last step used that all representables are finitely presented. The converse direction is easy.  $\square$

*Remark 4.4.* This style of proof crucially relies on the fact that  $\eta_{\mathcal{A}} : \mathcal{A} \rightarrow \text{Ind}(\mathcal{A})$  is fully faithful. This suggests that this style of argument fits into the setting of **Yoneda structures** [Wal17].

**Example 4.5.** For the monad  $\text{ldl}$  from theorem 3.6, given a  $X \in \text{Dcpo}$  an element  $x : 1 \rightarrow X$  is  $\text{ldl}$ -compact exactly when it is **compact or finite** in the sense of [Gie+80].

*Remark 4.6.* For the complete story to work one would need to work in the general setting of a virtual equipment to recover the right notion of generator, we will leave this for future work.

The following definition is a generalization of coherent morphisms of coherent locales [Joh82].

**Definition 4.7.** An algebra morphism  $f : (X, \alpha) \rightarrow (Y, \beta)$  is **coherent** if it preserves compact objects, that is if  $x : A \rightarrow X$  is compact then  $fx : A \rightarrow Y$  is compact.

**Proposition 4.8.** *Given algebras  $(X, \alpha)$  and  $(Y, \beta)$ , any left adjoint  $l : (X, \alpha) \rightarrow (Y, \beta)$  in the 2-category of algebras  $\mathbf{T}\text{-Alg}$  is coherent.*

*Proof.* Let  $r : (Y, \beta) \rightarrow (X, \alpha)$  be an algebra morphism which is right adjoint to  $l$  in  $\mathbf{T}\text{-Alg}$ . So it is in particular a right adjoint in  $\mathcal{B}$ .

We have to show that given a compact object  $x : A \rightarrow X$  of  $(X, \alpha)$ , we have that  $lx : A \rightarrow Y$  is compact in  $(Y, \alpha)$ . So take any diagram  $u : A \rightarrow \mathbf{T}Y$  in  $Y$  and notice

---

<sup>1</sup>Recall that an object  $X \in \mathcal{A}$  is finitely presented if for any filtered diagram  $D : I \rightarrow \mathcal{A}$

$$\mathcal{A}(X, \text{colim}_{i \in I} D(i)) \cong \text{colim}_{i \in I} \mathcal{A}(X, D(i))$$

$$\begin{array}{c}
\frac{\eta_Y l x \Rightarrow u}{T l \eta_X x \Rightarrow u} \\
\frac{\eta_X x \Rightarrow T r u}{x \Rightarrow \alpha T r u} \\
\frac{x \Rightarrow r \beta u}{l x \Rightarrow \beta u}
\end{array}$$

Thus  $l x$  is compact. □

## 5 Factorization

Next we introduce continuous algebras for lax-idempotent monads on a 2-category and describe a way of deducing from every lax-idempotent monad on a 2-category with enough limits an adjunction between the base category and the category of continuous algebras. A similar idea has been studied before in the descent theory for categories of modules by [Mes06], in which the conceptual meaning of the induced adjoint was not explored. We will show that the right adjoint from continuous algebras to the base 2-category takes a continuous algebras to its object of generators.

For the rest of this section we consider a fixed 2-category  $\mathcal{B}$  together with a lax-idempotent monad  $(T : \mathcal{B} \rightarrow \mathcal{B}, \mu, \eta)$ .

### 5.1 Continuous Algebras of Lax-Idempotent Monads

**Definition 5.1.** A **continuous algebra** for  $T$  is a triple  $(X, \alpha, \lambda_X)$  such that  $(X, \alpha)$  is a  $T$ -algebra and  $\alpha$  has a further left adjoint  $\lambda_X : X \rightarrow TX$ . A **morphism of continuous algebras** is a morphism of algebras  $f : (X, \alpha, \lambda_X) \rightarrow (Y, \beta, \lambda_Y)$  such that the mate  $\hat{f} : T f \lambda_X \cong \lambda_Y f$  of  $\tilde{f}$ , which we defined in theorem 3.9 is invertible

$$\begin{array}{ccccc}
X & \xrightarrow{\lambda_X} & TX & \xrightarrow{Tf} & TY \\
& \searrow \text{unit} & \downarrow \alpha & \nearrow \tilde{f}^{-1} & \downarrow \beta \\
& & X & \xrightarrow{f} & Y \\
& \searrow \text{id}_X & & & \searrow \lambda_Y \\
& & & & TY
\end{array}$$

$\text{id}_{TY}$  (curved arrow from  $TY$  to  $TY$ )  
 $\text{counit}$  (double arrow from  $TY$  to  $Y$ )

Notice that the triple of adjoints for a continuous algebra  $(X, \alpha, \lambda_X)$  always induces a 2-cell, which we will denote  $\theta_X : \lambda_X \Rightarrow \eta_X$ .

**Example 5.2.** For theorem 3.5, the continuous algebras are the completely distributive categories  $\mathbf{CDistCat}$  from [MRW12]. For theorem 3.6, the continuous algebras are the continuous posets [Gie+80], hence the name.

*Remark 5.3.* Every free algebra of a lax-idempotent monad is continuous by definition.

The next lemma should be taken as a categorification of: an element is compact in the sense of domain theory if and only if it is way below itself, see for example chapter 1 of [Gie+80].

**Lemma 5.4.** *For a continuous algebra  $(X, \alpha)$ , a generalized element  $x : A \rightarrow X$  is  $T$ -compact if and only if  $\theta_X x$  is invertible.*

*Proof.* ( $\Rightarrow$ ) Since we have that  $\eta_X x$  and  $\lambda_X x$  are both left lifts there is a canonical iso between them and it must invert  $\theta_X$ .

( $\Leftarrow$ ) Since  $\theta_X x$  is invertible,  $\eta_X x \cong \lambda_X x$ , and  $\lambda$  is a left adjoint and thus an absolute left lift of the identity, hence  $\eta_X x$  is an absolute left lift.  $\square$

**Corollary 5.5.** *The unit  $\eta_X : X \rightarrow TX$  is always compact.*

*Proof.* By theorem 5.3, any  $TX$  is a continuous algebra with the adjoint cylinder given by

$$\begin{array}{ccc} & T^2X & \\ \nearrow T\eta_X & \downarrow \mu_X & \nwarrow \eta_{TX} \\ TX & & TX \end{array}$$

Notice that naturality tells us that  $\eta_{TX}\eta_X \cong T\eta_X\eta_X$  and the fact that this isomorphism is given by  $\theta_{TX}$  can be found in [Koc95]. Thus the unit is compact.  $\square$

**Lemma 5.6.** *Coherent algebra morphisms between free algebras are automatically morphisms of continuous algebras.*

*Proof.* Take objects  $X, Y \in \mathcal{B}$  and an algebra morphism  $f : TX \rightarrow TY$  such that for any  $x : A \rightarrow X$  compact,  $fx : A \rightarrow Y$  compact. We want to show that  $Tf \circ T\eta_X \cong T\eta_Y f$ . By theorem 5.5 the unit  $\eta_X$  is always compact. Thus  $f\eta_X$  is compact by assumption. Thus by theorem 5.4:

$$\begin{aligned} TfT\eta_X\eta_X &\cong Tf\eta_{TX}\eta_X \\ &\cong \eta_{TY}f\eta_X \\ &\cong T\eta_Y f\eta_X \end{aligned}$$

Now notice that all the morphisms involved here are in fact algebra morphisms so by the universal property of free algebras

$$Tf \circ T\eta_X \cong T\eta_Y f.$$

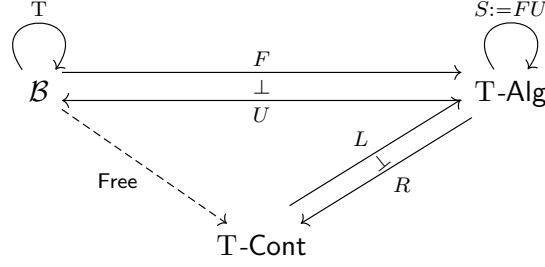
$\square$

*Remark 5.7.* In light of theorem 5.6 one could hope that coherent algebra morphisms are automatically in the image of *Free*. This is not the case. We will see a counterexample in Section 7.

**Lemma 5.8.** [Koc95] *The category of continuous algebras  $T\text{-Cont}$  for  $T$  is exactly the category of coalgebras for the comonad  $S := UF$  on algebras.*

$$\begin{array}{ccc} \begin{array}{c} T \\ \downarrow \\ \mathcal{B} \end{array} & \xrightleftharpoons[\quad U \quad]{\quad F \quad} & \begin{array}{c} S := FU \\ \downarrow \\ T\text{-Alg} \end{array} \end{array}$$

**Corollary 5.9.** *There is a comparison functor  $\text{Free} : \mathcal{B} \rightarrow T\text{-Cont}$  induced by the adjunction  $F \dashv U$ .*



*Free* is given on objects and morphisms by  $T$ . In particular this means that for any  $f : X \rightarrow Y$ ,  $Tf : TX \rightarrow TY$  is a morphism of continuous algebras.

**Lemma 5.10.** *Continuous algebra morphisms are coherent.*

*Proof.* Take a morphism of continuous algebras  $f : (X, \alpha, \lambda_X) \rightarrow (Y, \beta, \lambda_Y)$  and a compact object  $x : A \rightarrow X$  of  $X$ . And we would like to show that  $fx : A \rightarrow Y$  is compact in  $Y$ :

$$\frac{\frac{\frac{\eta_Y fx \Rightarrow u}{Tf \eta_X x \Rightarrow u}}{Tf \lambda_X x \Rightarrow u}}{\lambda_Y fx \Rightarrow u} \quad \frac{\lambda_Y fx \Rightarrow u}{fx \Rightarrow \beta u}$$

Which concludes the proof. □

## 5.2 A Factorization Theorem

**Definition 5.11.** An **adjoint cylinder** in a bicategory  $\mathcal{B}$  consists of two adjunctions  $l \dashv a \dashv u$  such that  $au \cong id$  and  $al \cong id$ . This configuration always induces a two-cell  $\theta : l \Rightarrow u$ .

**Example 5.12.** Every continuous algebra  $(X, \alpha, \eta_X)$  induces an adjoint cylinder  $\lambda_X \dashv \alpha \dashv \eta_X$ .

**Definition 5.13.** We say that a 2-category  $\mathcal{B}$  has **inverters of adjoint cylinders** if the induced two cell of any adjoint cylinder admits an inverter.

This can be seen as a lax version of having coreflexive equalizers.

**Proposition 5.14.** *If  $\mathcal{B}$  has inverters of adjoint cylinders, then there is a pseudofunctor  $K : T\text{-Cont} \rightarrow \mathcal{B}$  taking a continuous algebra to the inverter of its induced adjoint cylinder (cf. theorem 5.12).*

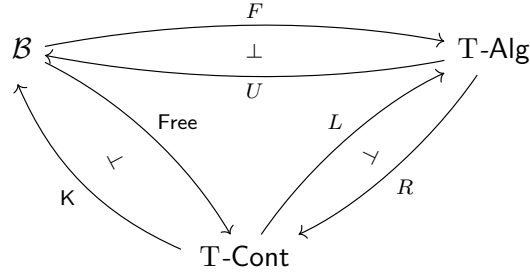
*Proof.* On objects  $K(X, \alpha, \lambda_X)$  is the inverter of  $\theta_X$ . By Section 5.1, any continuous algebra morphism  $f : (X, \alpha, \lambda_X) \rightarrow (Y, \beta, \lambda_Y)$  preserves compact objects, which by theorem 5.4 inverts the canonical 2-cell  $\theta_X : \lambda_X \Rightarrow \eta_X$  and thus induces a morphism between the inverters.

$$\begin{array}{ccc} \begin{array}{c} TX \\ \left( \begin{array}{c} \curvearrowright \\ = \theta_X \triangleright \\ \curvearrowright \end{array} \right) \\ X \end{array} & \xrightarrow{f} & \begin{array}{c} TY \\ \left( \begin{array}{c} \curvearrowright \\ = \theta_Y \triangleright \\ \curvearrowright \end{array} \right) \\ Y \end{array} \\ \uparrow \iota_X & & \uparrow \iota_Y \\ KX & \xrightarrow{\quad Kf \quad} & KY \end{array}$$

By Section 2.2 this assignment is functorial and that it extends to a pseudofunctor.  $\square$

Now we are ready to state the main theorem of this section.

**Theorem 5.15.** *If  $\mathcal{B}$  has inverters of adjoint cylinders and  $(T, \mu, \eta)$  is a lax-idempotent monad on  $\mathcal{B}$ , then  $\text{Free} \dashv K$  form a biadjunction between  $\mathcal{B}$  and the category of continuous algebras  $T\text{-Cont}$  making the left adjoints in the following diagram commute:*



*Proof.* The existence of the pseudofunctor  $K$  is by theorem 5.14. Explicitly, given a continuous algebra

$$\lambda \left( \begin{array}{c} TX \\ \downarrow \alpha \\ X \end{array} \right) \eta_X .$$

We can consider the inverter of the induced 2-cell  $\theta : \lambda \Rightarrow \eta_X$ :

$$\lambda \left( \begin{array}{c} TX \\ \xRightarrow{\theta} \\ X \end{array} \right) \eta_X$$

$$\begin{array}{c} X \\ \uparrow i_X \\ KX \end{array}$$

Which we denote by  $(KX, \iota_X)$ . The unit  $u_X : X \rightarrow KTX$  of the adjunction  $\text{Free} \dashv K$  is given by the universal property of  $\iota_X$ .

$$\begin{array}{ccc} & T^2X & \\ T\eta_X \left( \begin{array}{c} \xRightarrow{\theta} \\ \end{array} \right) \eta_{TX} & & \\ & TX & \\ \eta_X \nearrow & \uparrow i_{TX} & \\ X & \xrightarrow{u_X} & KTX \end{array}$$

The counit  $c_X : TKX \rightarrow X$  is given by the following composite:

$$TKX \xrightarrow{Ti_X} TX \xrightarrow{\alpha} X .$$



Notice that  $c_X$  is an algebra morphism.

The triangle identities will be essentially automatic once we show that  $c_X$  is indeed a morphism of continuous algebras.

Consider:

$$\begin{aligned} Tc_X T\eta_{KX} &\cong T(\alpha T\iota_X) T\eta_{KX} \\ &\cong T(\alpha T\iota_X \eta_{KX}) \\ &\cong T(\alpha \eta_X \iota_X) \\ &\cong T(\iota_X) \end{aligned}$$

and on the other hand:

$$\lambda c_X \cong \lambda \alpha T\iota_X$$

Now by Section 3.3,  $\lambda_X$  is an algebra morphism. Hence  $\lambda c_X$  and  $T(\iota_X)$  are algebra morphisms defined on the free algebra  $TKX$ . So we can use the universal property of free algebras to check isomorphism on the generators by precomposing with  $\eta_{KX}$ .

$$\begin{aligned} \lambda c_X \eta_{KX} &\cong \lambda \alpha T\iota_X \eta_{KX} \\ &\cong \lambda \alpha \eta_X \iota_X \\ &\cong \lambda i_X \\ &\cong \eta_X i_X \\ &\cong T\iota_X \eta_{KX} \end{aligned}$$

Thus  $\lambda c_X \cong Tc_X T\eta_{KX}$  making  $c_X$  is a continuous algebra morphism.

Therefore  $\tilde{F} \dashv K$ . □

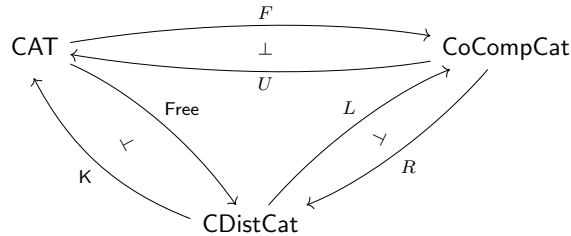
**Corollary 5.16.** *For any  $X \in \mathcal{B}$  admitting a T-algebra structure, we have that  $KTX \cong X$ .*

*Proof.* Since  $L\tilde{F} = F$  we have that  $KR \cong U$  by uniqueness of adjoints up to iso. Thus

$$\begin{aligned} KR(X, \alpha) &\cong K(TX, \mu_X) \\ &\cong U(X, \alpha) \\ &= X \end{aligned}$$

□

**Example 5.17.** For  $\mathcal{P} : \mathbf{CAT} \rightarrow \mathbf{CAT}$ , theorem 5.15 gives



The adjunction  $\text{Free} \dashv K$  is idempotent and the induced monad on  $\mathbf{CAT}$  is the usual Cauchy completion monad.

**Definition 5.18.** Going off the example  $\mathcal{P}$ , we call the induced monad  $K\text{Free} : \mathcal{B} \rightarrow \mathcal{B}$  the **T-Cauchy monad**.

We will see in Section 7 that this monad need not be idempotent.

### 5.3 Mnemonic Monads

In light of theorems 5.15 and 5.18 we would like to know when  $\mathbf{Free} \dashv \mathbf{K}$  becomes idempotent.

**Definition 5.19.** A lax-idempotent monad  $T : \mathcal{B} \rightarrow \mathcal{B}$  on a 2-category is **pre-mnemonic** if the inverter  $\iota_X : KTX \rightarrow TX$  of  $\theta_{TX} : T\eta_X \Rightarrow \eta_{TX}$  exists and the arrow  $Tu_X$  in the diagram

$$\begin{array}{ccc}
 T^2X & \xrightarrow{\quad} & T^2KTX \\
 \left( \begin{array}{c} \uparrow \\ \downarrow \end{array} \right) \begin{array}{c} \text{\scriptsize } T\eta_X \\ \text{\scriptsize } \dashv \mu_X \dashv \\ \text{\scriptsize } \eta_{TX} \end{array} & & \left( \begin{array}{c} \uparrow \\ \downarrow \end{array} \right) \begin{array}{c} \text{\scriptsize } \dashv \\ \text{\scriptsize } \dashv \end{array} \\
 TX & \xrightarrow{\quad Tu_X \quad} & TKTX \\
 \uparrow & \nwarrow \iota_X & \uparrow \\
 \eta_X & & \eta_{KTX} \\
 X & \xrightarrow{\quad u_X \quad} & KTX
 \end{array}$$

is an equivalence, where  $u_X$  is the canonical arrow into the inverter.

*Remark 5.20.* For pre-mnemonic monads the  $T$ -Cauchy monad is idempotent, in which case it truly is a kind of Cauchy completion.

**Definition 5.21.** A lax-idempotent monad  $T : \mathcal{B} \rightarrow \mathcal{B}$  on a 2-category is **mnemonic** if  $\eta_X$  is an inverter of  $\theta_{TX}$  :

$$\begin{array}{ccc}
 & T^2X & \\
 \left( \begin{array}{c} \uparrow \\ \downarrow \end{array} \right) \begin{array}{c} \text{\scriptsize } T\eta_X \\ \text{\scriptsize } \xrightarrow{\theta_{TX}} \\ \text{\scriptsize } \eta_{TX} \end{array} & & \\
 & TX & \\
 \uparrow & & \\
 \eta_X & & \\
 X & & 
 \end{array}$$

*Remark 5.22.* Mnemonic monads are a lax-idempotent variation of the monads of **descent type** as defined in [Mes06].

It is now easy to check that each mnemonic monad is in fact pre-mnemonic. Moreover we get the following easy consequences.

**Proposition 5.23.** *A lax-idempotent monad  $T : \mathcal{B} \rightarrow \mathcal{B}$  on a 2-category  $\mathcal{B}$  with inverters of adjoint cylinders is pre-mnemonic if and only if the induced adjunction  $\mathbf{Free} \dashv \mathbf{K}$  is idempotent. Moreover we call fixed points the **Cauchy complete objects** corresponding to  $T$  and the objects generated by their  $T$ -compact objects preserving the generators.*

**Corollary 5.24.** *A lax-idempotent monad  $T$  on a category with inverters of adjoint cylinders is mnemonic if and only if the induced functor  $\mathbf{Free}$  is a local equivalence.*

*Remark 5.25.* This holds even without the inverters existing.

**Theorem 5.26.** *Any pre-mnemonic monad  $T$  factors as an idempotent monad followed by a mnemonic monad.*

*Remark 5.27.* Mnemonic monads are always locally fully faithful in the sense of [Wal17], thus mnemonic monads are particularly close to Yoneda structures.

*Remark 5.28.* This means that a lax-idempotent monad is mnemetic if the base 2-category has Cauchy complete objects.

**Definition 5.29.** A continuous algebra  $(X, \alpha)$  has **enough T-compact objects** if the counit  $c_X : \mathrm{TK}X \rightarrow X$  is an equivalence.

From this we recover a similar theorem to that of [Mes06] for **nuclear adjunctions** [PH21] (adjunctions which are monadic and comonadic).

**Corollary 5.30.** *The left adjoint  $F : \mathrm{T}\text{-}\mathcal{A}lg \rightarrow \mathcal{B}$  is comonadic if and only if  $\mathrm{T}$  is mnemetic and each continuous algebra has enough compact objects.*

We are due for some examples.

**Example 5.31.** If we instantiate theorem 5.15 with the *Ind* construction on locally small categories we obtain an adjunction:

$$\begin{array}{ccc}
 & \mathrm{IndCat} & \\
 \swarrow \scriptstyle \perp & & \searrow \scriptstyle \top \\
 \mathrm{Cat} & \xrightleftharpoons[\mathrm{K}]{\mathrm{Ind}} & \mathrm{ContCat}
 \end{array}$$

Moreover since taking the compact elements of the ind-completion recover the Cauchy completion we have an induced equivalence, once we restricted to small categories

$$\mathrm{Cat}_{\mathrm{small}, \mathrm{cc}} \simeq \mathrm{Acc}_{\mathrm{fin}},$$

where  $\mathrm{Acc}_{\mathrm{fin}}$  has finitely accessible categories with functors preserving compact objects.

Moreover if we do the same for the *Ind*-completion on finitely cocomplete categories we get a kind of covariant Gabriel Ulmer Duality:

$$\mathrm{Rex} \simeq \mathrm{LFP}.$$

A usecase of the mnemeticity condition is that detecting free subalgebras simplifies. The proof can be found in the appendix C.6.

**Proposition 5.32.** *If  $\mathrm{T} : \mathcal{B} \rightarrow \mathcal{B}$  is mnemetic and  $\mathcal{B}$  has inverters of adjoint cylinders, then if  $A \in \mathcal{B}$  and  $(X, \alpha, \lambda_X)$  such that*

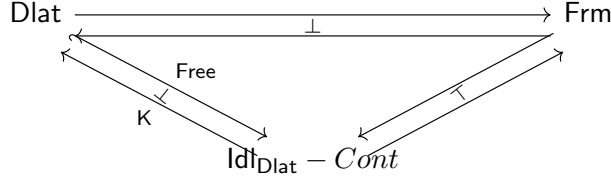
$$\begin{array}{ccc}
 X & \xrightleftharpoons[\scriptstyle l]{\scriptstyle i} & \mathrm{TA}
 \end{array}$$

*with  $i : X \rightarrow \mathrm{TA}$  a representably full and faithful continuous algebra morphism. In other words  $X$  is a reflective continuous subalgebra of  $\mathrm{TA}$ . Then  $c_X : \mathrm{TK}X \rightarrow X$  is an equivalence of continuous algebras.*

The following came out of conversations with Nathanael Arkor.

**Corollary 5.33.** *If  $(X, \alpha, \lambda)$  is a continuous algebra such that  $\lambda$  has a further left adjoint, then  $X$  is a free algebra. Thus the adjunction  $\mathrm{Cofree} - \mathrm{Forget}$  is nuclear.*

**Example 5.34.** We have that  $\mathbf{Frm} \rightarrow \mathbf{Dlat}$  is monadic with the left adjoint given on objects by  $\mathbf{Idl}$ . From this it is easy to see that the induced monad on  $\mathbf{Dlat}$  is mnemetic. Thus theorems 5.15 and 5.24 apply



The continuous algebras are continuous locales with a stronger condition on the generators, since  $\lambda : L \rightarrow \mathbf{Idl}_{\mathbf{Dlat}}(L)$  has to be a  $\mathbf{Dlat}$  morphism. These are called **stably locally compact locales** [Tow22]

Note that compactness stays the same as for  $\mathbf{Idl}$  on  $\mathbf{Pos}$ .

Now the essential image of  $\mathbf{Free}$  is exactly  $\mathbf{CohLoc}^{op}$  from [Joh82] giving us the known equivalence  $\mathbf{Dlat} \simeq \mathbf{CohLoc}^{op}$ .

*Remark 5.35.* There ought to be a similar version for toposes of this.

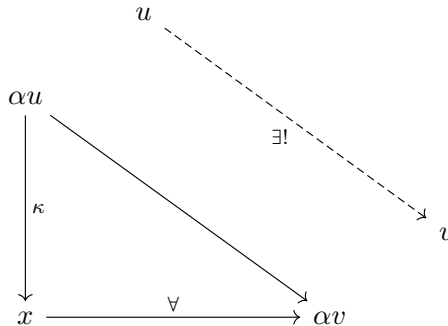
## 6 Way Below arrows and Slim Cocones

We now give a generic way of obtaining a way below relation from a lax-idempotent monad. Answering the folklorish question of how to generalize the wavy arrows appearing in the context of continuous categories.

**Definition 6.1.** Let  $T : \mathcal{B} \rightarrow \mathcal{B}$  be lax-idempotent, and let  $(X, \alpha)$  be an algebra.

For  $u : A \rightarrow TX$  and  $x : A \rightarrow X$ , we view 2-cells  $\kappa : \alpha \circ u \rightarrow x$  as (parametrized) **cocones**. Since in ordinary category theory, cocones are in correspondence with arrows out of colimits.

We call the cocone  $\kappa$  **slim**, if for every  $v : A \rightarrow TX$  and  $\phi : x \rightarrow \alpha \circ v$  there exists a unique  $\gamma : u \rightarrow v$  with  $\alpha * \gamma = \phi \kappa$ :



Moreover this property should be **absolute** in the sense that slim cocones are stable under right whiskering.

Slim cocones are in a sense absolute left lifting diagrams which have a leg pointing the wrong way.

We say that  $\kappa$  is **colimiting** if it is invertible.

**Proposition 6.2.** *If  $\kappa : \alpha u \Rightarrow x$  is a slim colimiting cocone, then  $(u, \kappa^{-1})$  is an absolute left lift:  $u = \mathbf{lift}_{\alpha} x$  with an invertible cell.*

**Corollary 6.3.** *Slim colimiting cocones are unique up to isomorphism.*

**Corollary 6.4.**  *$A$  is continuous iff every parametrized object  $a : X \rightarrow A$  is the vertex of a slim colimit cocone  $\alpha u \rightarrow a$ .*

*Proof.* ( $\Rightarrow$ ) Since  $A$  is continuous we have a furthest left adjoint  $\lambda \dashv \alpha$  which is always an absolute left lift  $\text{llift}_\alpha id_A$ , moreover since it is a coreflection the two cell is invertible. Thus precomposing with any parametrized object  $a : X \rightarrow A$  gives an absolute left lift with invertible two cell  $\lambda a = \text{llift}_\alpha a$ .

( $\Leftarrow$ ) Take the slim colimit cocone associated to  $id_A : A \rightarrow A$  to obtain an absolute  $\text{llift}_\alpha id_A$  together with an invertible two cell, thus this left lift will be the desired left adjoint.  $\square$

**Definition 6.5.** Let  $T : \mathcal{B} \rightarrow \mathcal{B}$  be lax-idempotent, and let  $(X, \alpha)$  be an algebra.

A **way-below arrow** between generalized elements  $x, y : A \rightarrow X$  is a 2-cell  $\gamma : x \rightarrow y$  such that the 2-cell obtained from the counit of the adjunction defining the algebra structure  $\alpha \eta_X x \rightarrow y$  is a slim cocone.

**Proposition 6.6.** *A parametrized object  $x : A \rightarrow X$  is compact iff the identity  $x \rightarrow x$  is a way-below arrow*

## 7 The counterexample

**Definition 7.1.** A subset  $C \subseteq D$  of a dcpo  $D$  is **Scott-closed** if it is down-closed and closed under directed joins.

**Lemma 7.2.** (Cf. Lemma 4 [VT04]) *There is a left adjoint  $\text{cl} : \text{Dcpo} \rightarrow \text{Sup}$  making the forgetful functor  $U : \text{Sup} \rightarrow \text{Dcpo}$  monadic. It sends a dcpo  $D$  to its lattice  $\text{cl}(D)$  of Scott-closed sets. Moreover the induced monad  $\text{cl} := U \circ \text{cl}$  is lax-idempotent.*

*Remark 7.3.* The lax idempotence can be seen from the framework of conservative cocompletions [Lam]. Since  $X \in \text{Dcpo}$ ,  $\text{cl}(X)$  adds in the *missing* finite joins while keeping the directed joins.

The following proposition tells us that we indeed recover the right notion of compactness for  $\text{cl}$  from [HZ09], which they called *C-compactness*.

**Proposition 7.4.** *An element  $x : 1 \rightarrow L$  of a suplattice is **cl-compact** if and only if*

$$c \leq \bigvee A \Rightarrow c \in A$$

for all  $A \in \text{cl}(L)$ .

*Proof.* By definition  $x$  is *cl-compact* iff  $\downarrow x = \text{lift}_x \bigvee$ , that is

$$\downarrow x \subseteq_c lA \Leftrightarrow x \leq \bigvee A.$$

But the left to right direction is trivial giving us the desired result.  $\square$

**Proposition 7.5.** *Let  $D$  be a dcpo, and  $\mathcal{A} \subseteq_{\text{cl}} \text{cl}(D)$  a Scott-closed set of Scott closed sets. Then  $\bigcup \mathcal{A}$  is Scott-closed in  $D$ .*

*Proof.* Consider a directed family  $a : (I, \leq) \rightarrow \bigcup \mathcal{A}$ . Then for every  $i \in I$  there exists an  $A_i \in \mathcal{A}$  with  $a_i \in A_i$ . Define  $A'_i = \bigcap_{j \geq i} A_j$ . Then the  $A'_i$  are also all closed and in  $\mathcal{A}$  and contain  $a_i$ . We have  $a = \bigvee_i a_i \in \bigcup_i A'_i \subseteq \bigvee_i A'_i \in \mathcal{A}$  and therefore  $a \in \bigcup \mathcal{A}$ .  $\square$

*Remark 7.6.* The equalizer (inverter)

$$\mathbf{K}(TD) \rightarrow TD \rightrightarrows T^2D$$

consists precisely of the cl-compact elements of  $TD$ . It clearly contains all principal downsets, but there there might be more!

Consider Johnstone's dcpo  $\mathcal{J}$  [Joh06] whose underlying set is  $\mathbb{N} \times \mathbb{N}_\infty$  and where  $(i, n) \leq (j, m)$  iff either  $i = j$  and  $n \leq m$  or  $m = \infty$  and  $n \leq j$ .

**Proposition 7.7.**  $\mathcal{J}$  is cl-compact in  $\text{cl}(\mathcal{J})$ .

*Proof.* Assume that  $\mathcal{J}$  is the supremum (and therefore the union) of a closed set  $\mathcal{A} \subseteq \text{cl}(\mathcal{J})$ . Then for all  $i$  there is an  $A_i \in \mathcal{A}$  such that  $(i, \infty) \in A_i$  and therefore  $B_i = \mathbb{N} \times \{0, 1, \dots, i\} \in \mathcal{A}$  since it's a closed subset of  $A_i$ . Then  $J$  is the directed join of the  $B_i$  and it must be contained in  $\mathcal{A}$  since  $\mathcal{A}$  is assumed to be closed.  $\square$

We observe that  $T\mathcal{J}$  is  $\mathcal{J}$  with a new greatest element adjoined. But in  $\mathbf{KT}\mathcal{J}$ , the original  $\mathcal{J}$  is again closed and cl-compact and not principal, so  $(\mathbf{KT})^2\mathcal{J}$  adds yet another element *under* the greatest element added in the first step. It seems that transfinite iteration never stops.

**Corollary 7.8.** *The induced adjunction  $\mathbf{Free} \dashv \mathbf{K}$  from theorem 5.15 need not be idempotent. That is there are non pre-mnematic lax-idempotent monads.*

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## A Background on Pseudomonads

Pseudomonads encode structure which is defined up to isomorphism, for example a monoidal structure on a category. A complete set of references for the constructions in this section is [Lei98; BKP89; Lac00].

**Definition A.1.** Let  $\mathcal{B}$  and  $\mathcal{C}$  be 2-categories, a **pseudofunctor**  $F : \mathcal{B} \rightarrow \mathcal{C}$  consists of the following data:

- A function  $F_0 : \mathcal{B}_0 \rightarrow \mathcal{C}_0$ , where for  $A \in \mathcal{B}$  we denote  $F_0(A)$  as  $F(A)$
- For each pair of objects  $A, B \in \mathcal{B}$ , a functor  $F_{A,B} : \mathcal{B}(A, B) \rightarrow \mathcal{C}(A, B)$  where we again denote  $F_{A,B}(f)$  as  $F(f)$
- For each object  $A \in \mathcal{B}$  an invertible 2-cell  $F_{\text{id}_A} : \text{id}_{F(A)} \Rightarrow F(\text{id}_A)$
- For each triple  $A, B, C \in \mathcal{B}$  a natural isomorphism (natural in  $f : A \rightarrow B$  and  $g : B \rightarrow C$ )  $F_{g,f} : F(g) \circ F(f) \Rightarrow F(g \circ f)$

This datum is required to make two diagrams commute which appear in [Lei98].

**Definition A.2.** Given 2-categories  $\mathcal{B}$  and  $\mathcal{C}$  with pseudofunctors  $F, G : \mathcal{B} \rightarrow \mathcal{C}$  between them, a **pseudonatural transformation**  $\alpha : F \Rightarrow G$  is given by:

- For each object  $A \in \mathcal{B}$ , an arrow  $\alpha_A : F(A) \rightarrow G(A)$  in  $\mathcal{C}$
- For each arrow  $f : A \rightarrow B \in \mathcal{B}$ , an invertible 2-cell  $\alpha_f : G(f) \circ \alpha_A \Rightarrow \alpha_B \circ F(f)$ .

And this datum has to satisfy coherence diagrams [Lei98].

**Definition A.3.** A **modification** between pseudonatural transformations  $\alpha, \beta : F \Rightarrow G$ , denoted  $\theta : \alpha \Rightarrow \beta$  is given by a family of 2-cells  $\theta_A : \alpha_A \Rightarrow \beta_A$  indexed by objects  $A \in \mathcal{B}$  such that for each arrow  $f : A \rightarrow B$  the following square commutes:

$$\begin{array}{ccc} G(f)\alpha_A & \xrightarrow{G(f)\theta_A} & G(f)\beta_A \\ \alpha_f \downarrow & & \downarrow \beta_f \\ \alpha_B F(f) & \xrightarrow{\theta_B F(f)} & \beta_B F(f) \end{array} .$$

**Definition A.4.** Let  $\mathcal{B}$  be a 2-category, a **pseudomonad**  $(T : \mathcal{B} \rightarrow \mathcal{B}, \mu, \eta)$  consists of the following data:

- A pseudofunctor  $T : \mathcal{B} \rightarrow \mathcal{B}$
- pseudonatural transformations:  $\mu : T^2 \Rightarrow T$ ,  $\eta : \text{id} \Rightarrow T$
- invertible modifications:  $\lambda : \mu \circ \eta T \Rightarrow \text{id}$ ,  $\rho : \mu \circ T \eta \Rightarrow \text{id}$  and  $\alpha : \mu \circ T \mu \Rightarrow \mu \circ \mu T$ .

These are required to satisfy two diagrams which can be found in [Lac00].

## B Noetherian objects

With the greatest amount of bad intent one could call an algebra  $(X, \alpha)$  of a lax-idempotent monad  $T$  on a 2-category **Noetherian** if  $\text{id}_X : X \rightarrow X$  is compact in the sense of Section 4.

It is not too hard to show that these are the fixed points of the monad.

**Lemma B.1.** *An algebra  $(X, \alpha)$  is Noetherian if and only if  $\alpha : TX \rightarrow X$  is an equivalence.*

*Proof.* The only if direction is clear. For the other direction, take  $(X, \alpha)$  Noetherian, so  $\text{id}_X$  is compact, thus  $\eta_X = \eta_X \text{id}_X = \text{lift}_\alpha \text{id}_X$  is an absolute left lift. Thus by Section 2.2

$$\eta_X \dashv \alpha.$$

Thus  $\eta_X$  is an algebra morphism by Section 3.3. Hence by the universal property of free algebras we have that

$$\eta_X \alpha \eta_X \cong \eta_X,$$

which implies that  $\eta_X \alpha \cong \text{id}_{TX}$ . Thus  $\alpha$  is an equivalence.  $\square$

**Proposition B.2.** *Any reflective subalgebra of a Noetherian algebra  $(X, \alpha)$  is Noetherian.*

*Proof.* Take a reflective subalgebra  $A \subseteq X$  with inclusion reflection  $r \dashv i$ . By Section 4,  $r$  preserves compact objects, thus by absoluteness of compact objects  $r \circ \text{id}_X \circ i$  is compact in  $A$ . Hence  $\text{id}_A$  is compact in  $A$ , so  $A$  is Noetherian.  $\square$

**Proposition B.3.** *Any algebra coreflection of a Noetherian algebra  $(X, \alpha)$  is Noetherian.*

*Proof.* Same proof.  $\square$

**Corollary B.4.** *Noetherian spaces are closed under quotients and subspaces.*

## C Definable compactness

We can do a tiny bit better than 5.15 even if we don't have inverters.

**Definition C.1.** Compactness is **definable** for an algebra  $(X, \alpha)$  if there is a universal compact arrow  $\iota : KX \rightarrow X$ , that is:

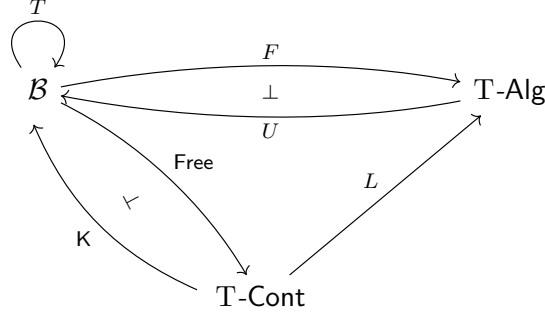
- $\iota_X$  is compact
- For every object  $Y \in \mathcal{B}$ :  $- \circ \iota : \mathcal{B}(Y, X) \rightarrow \mathcal{B}(Y, A)$  is full and faithful with its replete image in compact arrows for  $X$

This defines the 2-category  $T\text{-Alg}_{\text{def}}$  of algebras with definable compactness, coherent algebra morphisms and 2-cells of  $\mathcal{B}$ .

**Proposition C.2.** *Compactness is definable for continuous algebras precisely when the inverter of their canonical 2-cell exists.*

With this definition in hand we can drop the restrictions of theorem 5.15.

**Theorem C.3.** Any lax-idempotent monad  $T : \mathcal{B} \rightarrow \mathcal{B}$  induces an adjunction between  $\mathcal{B}$  and  $T\text{-Alg}_{\text{def}}$  such that  $L \circ \text{Free} \cong F$ :



**Definition C.4.** An algebra  $(X, \alpha)$  with definable compactness ‘has enough compact objects’ if the counit  $c_X : TKX \rightarrow X$  is an equivalence.

**Corollary C.5.** A lax-idempotent monad is mnemetic precisely if  $\text{Free}$  is a local equivalence and pre-mnemetic precisely if  $\text{Free} \dashv K$  is idempotent.

We now quickly return to the proof of theorem 5.32:

**Proposition C.6.** Assume  $T : \mathcal{B} \rightarrow \mathcal{B}$  is mnemetic and  $\mathcal{B}$  has inverters of adjoint cylinders. And assume  $A \in \mathcal{B}$  and that  $(X, \alpha, \lambda_X)$  is a continuous algebra such that

$$X \begin{array}{c} \xrightarrow{i} \\ \xleftarrow{\tau} \\ \xleftarrow{l} \end{array} TA$$

with  $i : X \rightarrow TA$  a representably full and faithful continuous algebra morphism. In other words  $X$  is a reflective continuous subalgebra of  $TA$ . Then  $X$  has enough compact objects, that is  $c_X : TKX \rightarrow X$  is an equivalence of continuous algebras.

*Proof.* We construct a pseudoinverse to  $c_X$ . Denote  $\iota_X : KX \rightarrow X$  the canonical inclusion and notice that by 5.5  $\eta_A$  is compact. By Section 4  $l$  preserves compact objects, thus  $l\eta_A : A \rightarrow X$  is compact. The universal property of  $\iota_X$  gives us a factorizing morphism  $f : A \rightarrow KX$  such that  $\iota_X f \cong l\eta_A$

$$\begin{array}{ccc} X & \xleftarrow{l} & TA \\ \uparrow \iota_X & & \uparrow \eta_A \\ KX & \xleftarrow{\quad f \quad} & A \end{array}$$

Let  $g : X \rightarrow TKX$  be the composite

$$X \xrightarrow{i} TA \xrightarrow{Tf} TKX .$$

Since  $i$  is a continuous algebra morphism,  $g$  is continuous algebra morphism.

So  $g \circ c_X : TKX \rightarrow TKX$  is a continuous algebra morphism defined on free algebras, thus by mnemeticity  $\eta_{KX} K(g \circ c_X) \cong g \circ c_X \circ \eta_{KX}$ . Now since we are working with algebra

morphisms defined over free algebras we can check equivalences on generators

$$\begin{aligned} g \circ c_X \eta_{\mathbf{K}X} &= T(f) i \alpha T(\iota_X) \eta_{\mathbf{K}X} \\ &\cong T(f) i \alpha \eta_X \iota_X \\ &\cong T(f) i \iota_X. \end{aligned}$$

Thus  $T(f) i \iota_X \cong \eta_{\mathbf{K}X} \mathbf{K}(g \circ c_X)$ . So precomposing with  $c_X$  gives us on one hand

$$c_X \eta_{\mathbf{K}X} \mathbf{K}(g \circ c_X) \cong \iota_X \mathbf{K}(g \circ c_X).$$

And on the other hand

$$\begin{aligned} c_X T(f) i \iota_X &= \alpha T(\iota_X) T(f) i \iota_X \\ &\cong \alpha T(\iota_X f) i \iota_X \\ &\cong \alpha T(l \eta_A) i \iota_X \\ &\cong \alpha T(l) T(\eta_A) i \iota_X \\ &\cong l \mu_A T(\eta_A) i \iota_X \\ &\cong l i \iota_X \\ &\cong \iota_X \end{aligned}$$

Hence  $\mathbf{K}(g \circ c_X) \cong \text{id}_{\mathbf{K}X}$  and by theorem 5.24

$$(g \circ c_X) \cong \text{Free}(\text{id}_{\mathbf{K}X}).$$

And the other direction is now immediate

$$\begin{aligned} c_X g &= \alpha T(\iota_X) T(f) i \\ &\cong \alpha T(\iota_X f) i \\ &\cong \alpha T(l \eta_A) i \\ &\cong \alpha T(l) T(\eta_A) i \\ &\cong l \mu_A T(\eta_A) i \\ &\cong l i \\ &\cong \text{id}_X \end{aligned}$$

So  $c_X$  is an equivalence of continuous algebras. □

*Remark C.7.* A similar argument shows that assuming  $T$  is lax idempotent and  $X \rightarrow TA$  is a reflective subalgebra, then  $c_X$  has a section. The problem is that  $c_X$  need not be representably fully faithful in general as in Section 7.

## D A big list of examples

In this final section we wanted to show that lax-idempotent monads come in many forms each giving rise to different interesting notions of generator!

The shape column indicates the kind of diagrams that we cocomplete under, all means all

(small) diagrams, finite means only finite diagrams, etc. The case of cocartesian arrows is in reference to turning an arbitrary functor  $F : E \rightarrow B$  over a fixed base  $B$  into an opfibration  $OP(F) : F \downarrow B \rightarrow B$ . This defines the action of a lax-idempotent monad [Koc13].

Question marks correspond to notions of compactness we don't have a name for yet.  $\mathcal{D}$  and  $\mathcal{D}_{fg}$  are the down-set constriction and the finitely generated down-set construction respectively. **Fam** is the families construction/ free coproduct completion. **Reg** is the reg/lex completion, **Ex** is the ex/lex completion. And finally **Env** is the envelope of a multicategory, which has been shown to be lax-idempotent in for example [Her01].

2-Category	Shape	L.I. Monad	Algebras	T-Compactness	Reference
Pos	all	$\mathcal{D}$	Sup	Completely-join-prime element	[GG24]
Pos	directed	Idl	Dcpo	join-prime element	[Gie+80]
Pos	finite	$\mathcal{D}_{fg}$	$\vee$ -Lat	finite/compact element	[GG24] (Exercise 1.3.8)
CAT	all	$\mathcal{P}$	CoCompCat	atomic object	[PT21]
CAT	directed	Ind	IndCoCompCat	finitely presented object	[Joh82] (Ch. VI)
CAT	discrete	Fam	CoprodCocompCat	connected object	[CJ95]
CAT	sifted	SInd	SIndCoCompCat	?	[AR01]
Lex	image factorizations	Reg	RegCat	regular projective object	[CV98]
Lex	effective quotients	Ex	ExCat	effective projective object	[CV98]
AddCat	effective quotients	Ex	AbCat	effective projective object	[RV01]
CAT/ $B$	cocartesian arrows	OP	OpFib( $B$ )	?	[Koc13]
MultiCat	representables	Env	MonCat	?	[Her01]

This last example  $Env : MultiCat \rightarrow MultiCat$  is special in another way. Namely it is an example of theorem 5.30, which has been shown in [EM07] for symmetric Multicategories.